



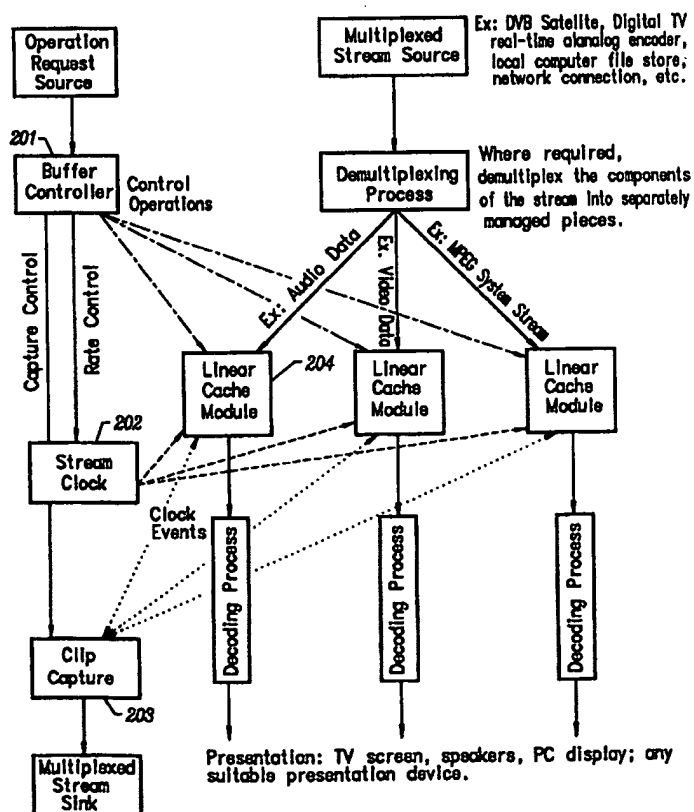
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(54) Title: METHOD AND APPARATUS IMPLEMENTING RANDOM ACCESS AND TIME-BASED FUNCTIONS ON A CONTINUOUS STREAM OF FORMATTED DIGITAL DATA

(57) Abstract

A continuous stream of formatted digital data, such as a video segment, audio segment, or information stream, appears to be a fixed length segment under certain circumstances, defining a virtual segment within the continuous stream which moves forward in time in synchrony with the continuous stream. The virtual segment thus defined can be explored in a non-linear fashion at arbitrary playback rates. For instance, concepts such as rewind, pause, frame advance, and fast forward become meaningful even though the continuous stream never ceases.



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Method And Apparatus Implementing Random Access And Time-Based Functions On A Continuous Stream Of Formatted Digital Data

BACKGROUND OF THE INVENTION

TECHNICAL FIELD

The invention relates to digital signal processing. More particularly, the invention relates to a family of methods which provide for the pass through or capture of linear streams of digital information represented in various formats.

DESCRIPTION OF THE PRIOR ART

With the advent of consumer audio and video products that employ sophisticated digital signal processing techniques it is becoming necessary to find ways to exploit the full potential of digital technology. For example, it would be desirable to provide methods for the pass through or capture of linear streams of digital information represented in various formats, while at the same time providing the appearance to the consumer of a locally stored digital stream that allows for the repositioning and playback of virtual segments of the apparently local digital stream.

A mechanical device which performs some of these functions is the Video Cassette Recorder (VCR), which uses a magnetic tape to store the information. The inherently linear nature of tape leads to functions such as rewind, fast forward, and pause. However, a VCR cannot both capture and play back information at the same time, so it cannot be used to implement this capability.

Linear streams of information are a fixture of modern life. Consider broadcast radio stations, broadcast television stations, satellite broadcasts, cable television, video tapes, and compact disks. Increasingly, such information is represented in a fashion suitable for manipulation by automated electronic hardware, such as computers or media decoders. For example, the Direct Video Broadcast (DVB) standards address digital broadcasting from satellites, terrestrial stations, and cable television systems. Even analog

2
broadcasts, such as normal NTSC (National Television Standards Committee) broadcasts from familiar local stations, may be captured and digitized in real time by modern equipment, making them appear to be linear digital streams.

Though such streams never terminate, and an individual viewer of the stream is unable to effect directly how such streams are delivered, it is desirable to provide the illusion for the consumer that recent portions of the stream are stored locally in some manner, such that typical VCR-like functions can be performed on the stream, *e.g.* pause, rewind, and fast forward. The desire for this capability arises from the fact that the schedule and timing of the broadcast almost never matches the needs of the individual viewer. For instance, the viewer may wish to stop the stream for a few moments to discipline an unruly child. Or perhaps the viewer's attention was distracted from the stream for a few moments, causing him to miss a critical scene, in which case the viewer would like to rewind to the point he missed and play it again.

Ideally, a device local to the viewer should capture the entire stream as it is being broadcast and store it in some manner. For example, if two video tape recorders are available, it might be possible to Ping-Pong between the two. In this case, the first recorder is started at the beginning of the program of interest. If the viewer wishes to rewind the broadcast, the second recorder begins recording, while the first recorder is halted, rewound to the appropriate place, and playback initiated. However, at least a third video tape recorder is required if the viewer wishes to fast forward to some point in time after the initial rewind was requested. In this case, the third recorder starts recording the broadcast stream while the second is halted and rewound to the appropriate position. Continuing this exercise, one can quickly see that the equipment becomes unwieldy, unreliable, expensive, and hard to operate, while never supporting all desired functions. In addition, tapes are of finite length, and may potentially end at inconvenient times, drastically lowering the value of the solution.

It is possible to implement this capability using a digital computer, where digital streams are stored in some fashion analogous to video tape and where the computer performs the switching between the various virtual tape decks. Even using a digital computer, this strategy suffers from the same weaknesses as the physical system above. It would be desirable to avoid

these issues by providing a technique for storing the streams of information on a temporary basis.

When using a digital computer to perform any technique which achieves this functionality, there are a number of issues which must be taken into account for proper operation. The first of these is storage of the broadcast stream. Within a digital computer, a stream of information is represented as a sequence of blocks of digital data. For example, when encoding an NTSC television broadcast stream, each field of analog data is converted to a block of 8-bit digital samples representing the field. If the analog signal is faithfully represented, each digital block contains approximately 0.5 MB of data, one second of video requires approximately 30 MB of storage, and 30 seconds of video requires approximately 900 MB of storage, greater than the capacity of a compact disc. Manipulation of video in this form clearly becomes unworkable when any useful length of stored video is contemplated.

As an example, consider U.S. patent no. 5,625,46, which concerns the use of a magneto-optic disk for the storage of broadcast television transmissions. The amount of storage available on such media is currently about 5 to 10 gigabytes, which is sufficient for approximately 5 seconds of video storage - clearly insufficient. In addition, the device disclosed does not permit the simultaneous recording and playback of the same program.

Limited storage capacity is dealt with by compressing the video stream using an algorithm, typically one of the MPEG (Moving Pictures Experts Group) standard algorithms, which can achieve a useful compression of 100:1 in many instances. MPEG video is represented as a sequence of Groups Of Pictures (GOPs), in which each GOP begins with an index frame, called the I-frame. The I-frame is a block of digital data which is compressed using Discrete Cosine Transform (DCT) and other techniques, similar to the still-picture Joint Photographic Experts Group (JPEG) standard.

The GOP may represent up to 15 additional frames by providing a much smaller block of digital data that indicates how small portions of the I-frame, referred to as macroblocks, move over time. Thus, MPEG achieves its compression by assuming that only small portions of an image change over time, making the representation of these additional frames extremely compact.

Unlike the uncompressed data example above, or examples based on video tape recording, each frame is thus represented as a variable length block of binary data. Additionally, although GOPs have no relationship between themselves, the frames within a GOP have a specific relationship which builds off the initial I-frame. Thus, any method which stores a digitized stream and allows random access to the stored information must take into account the variable (and unpredictable) data sizes involved, as well as be cognizant of the relationships between blocks of the stream.

A second issue for a digital computer based implementation of such methods is that multiple streams of information must be handled in parallel. For example, a broadcast stream is actually composed of at least two unique sequences of information, *i.e.* a stream of digital blocks representing the visual image and a stream of digital blocks representing the audible image. If the audio effect is instead stereo, then two audio streams are included, each unique. A broadcast signal may have additional data, such as the Secondary Audio Program (SAP), where the stream of information is a translation of the audio signal to a different language. Another stream which may be present is the Closed Caption (CC) stream, which provides a textual representation of spoken language in the audio stream(s). The simple broadcast stream described earlier may therefore have at least five different components, each one compressed using different techniques. When presenting this complex stream to a viewer, the blocks of each stream must be decoded at appropriate times for the compression methods involved and synchronized with the presentation of all other streams.

Also of interest are digital broadcasting technologies, such as DVB. A DVB channel is formed in an MPEG2 Transport Multiplex, which is an encoding scheme that provides for interleaving any number of discrete streams of digital information into a single stream of digital data, using techniques based on Time Division Multiplexing (TDM). The example television signal above can be encoded into a DVB channel using five discrete streams, leaving additional capacity for other streams.

There is increasing interest in adding additional information streams to a standard broadcast signal. For instance, it may be desirable to transmit audio channels in several different languages in parallel with the video stream. Or, perhaps information that is interpreted as a Web page is broadcast in such a way as to be synchronized with the video to provide a multimedia

presentation. The number of streams which must be synchronized may be arbitrary, and each stream may be represented using different and unique storage and compression techniques which have their own synchronization requirements and inter-frame relationships.

Any methods which provide functionality similar to that described above using some form of digital computer must contain techniques which resolve these issues.

SUMMARY OF THE INVENTION

The methods and apparatus described herein implement a novel and unique facility that provides for the manipulation of a continuous stream of digital information as if it supported random access and variable rate presentation on demand by the consumer of the stream. The possible operations on the data stream are limited only by the storage capacity and IO bandwidth of the system which implements the apparatus and methods, whether that storage be volatile, such as DRAM, or non-volatile, such as hard disk storage, or a combination of both. The apparatus and methods support operations, sometimes referred to as virtual VCR functions, such as pause, rewind, fast forward, and play, as well as more sophisticated and unique operations, such as play faster, play slower, and play in reverse. Additionally, the apparatus and methods described herein provide a mechanism whereby such operations are instantaneous from the point of view of the consumer of the data stream, unlike in mechanical systems. The apparatus and methods also support the ability to capture portions of the data stream which are temporarily held by the apparatus, such that this captured data can be stored or sent over a network in formats suitable for presenting the data to other consumers, or for editing and manipulation by other tools or methods.

The invention provides a method and apparatus for providing pass through or capture of continuous linear streams of digital information represented in various formats while providing the appearance of a locally stored stream. The preferred embodiment of the invention comprises at least one media cache for copying blocks of data from the information stream. Data in the media cache can be viewed as a snapshot of the continuous stream of digital information. The invention also comprises a playback pointer. The playback pointer position selects a portion of the media cache that is to be accessed to

provide functions including any of pause, rewind, fast forward, play, play faster, play slower, and play in reverse.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a schematic representation of a continuous media buffer according to the invention;

Fig. 2 is a block schematic diagram of a system incorporating a continuous media buffer according to the invention;

Fig. 3 is a timing diagram showing the operation of a linear cache according to the invention;

DETAILED DESCRIPTION OF THE INVENTION

The invention contemplates a new and unique apparatus for transforming and manipulating continuous streams of digital information using the methods and techniques described herein.

Figure 1 provides a general overview of how the preferred embodiment of the invention operates on a continuous stream of data. In particular, one or more media caches 10 copy blocks of data from an information stream 12. The position of a playback pointer 14 selects a portion of the media cache that is to be accessed, thus allowing such functions as pause, rewind, fast forward, and play, as well as more sophisticated and unique operations, such as play faster, play slower, and play in reverse.

The apparatus is independent of the format of the underlying digital data stream. To apply the invention, it is only necessary that the underlying data stream is composed of a number of discrete blocks of data, and that only a small set of attributes are associated with each block. The data blocks themselves may be of arbitrary size and format. It is assumed that processes specialized to a data source or to a decoding operation may proceed freely within the context of the invention, indirectly controlled by the methods described herein.

Figure 2 illustrates the preferred embodiment of this invention. While the example in Figure 2 is provided in connection with an MPEG stream of

information, it should be appreciated that the invention is readily applied to any digital stream of information.

Manipulation and transformation of a data stream is performed in the invention by a combination of the following elements:

1. The Buffer Controller (201) is the main management device. It accepts external requests (generated, for instance, from a remote control device) for operations on the digital stream and, in turn, generates appropriate control messages for the other devices of the invention;
2. The Stream Clock (202) provides a general device for synchronizing operations on a set of linear caches, such that multiple streams of data which must be presented in a synchronized fashion are correctly positioned, and that they present their data at consistent delivery rates;
3. The Linear Cache (204) is a general device for buffering the information contained in a stream of digital information, such that the data in the cache can be viewed as a snapshot of the continuous stream of digital data; and
4. The Clip Capture (203) device is a general mechanism for capturing all or part of a cached set of information streams and presenting the resulting data as an organized stream of data to a data sink such as, for example, a computer file system or another display device.

The following sections describe each of the foregoing components, and how they interact, in detail.

BUFFER CONTROLLER

The Buffer Controller (BC) module is responsible for instantiating the methods that provide for operation of the invention for any particular stream of data. The preferred embodiment of the invention includes one Stream Clock (SC) module, zero or more Clip Capture (CC) modules, and one or more Linear Cache (LC) stream modules. Operation of the invention preferably relies on designating one of the LC streams as the key stream. All positioning of playback from the LCs is based on first positioning the key stream, and then using the final position of that stream to position all other

streams properly. For example, if the key stream is an MPEG video stream, then positioning is based on aligning the current block pointer to the beginning of an MPEG Group Of Pictures (GOP). Once the key stream is positioned, the BC repositions all other streams to be synchronized with the key stream.

The BC is also responsible for controlling the clock rate of the Stream Clock, which operates as described below. The rate at which the SC delivers clock events to the Linear Caches is based on the rate specified by the BC. Thus, the BC can speed up or slow down the playback of the entire collection of streams by requesting the SC to adjust its clock rate.

Finally, the BC is responsible for initiating clip capture operations using the Clip Capture method.

STREAM CLOCK

The Stream Clock provides a central synchronization facility that distributes time-based events to a number of LCs. Each stream of information in a broadcast program is encoded using different techniques, each technique having unique notions of, for example, block size, data format, and presentation time stamps for decoding. For instance, in the time it takes for a single MPEG video frame to be decoded and presented, there may be several MPEG audio frames to be decoded and presented. The Stream Clock distributes events to each LC at the proper rate for that LC.

The Stream Clock uses a single queue for managing events to be distributed using standard, prior art techniques that are familiar to those skilled in the art of building time-based software systems.

Operation of the Stream Clock is detailed below.

Events are not queued with an absolute time, but with the relative offset from the previously queued event. The first event in the queue is marked with a time which is the offset from the current time. The Stream Clock then requests that the underlying operating system invoke a call out when that amount of time has passed, *e.g.* the SC queues a timer event. When the call out is invoked, the SC dispatches the first clock event and then each succeeding event whose offset is zero, *i.e.* each event that is to occur at the same time as the previous event.

Each event in the queue is also marked with a time interval and the absolute time of the data stream that the next event should occur. After a queued event is dispatched, the absolute time of the next event is composed by adding the time interval to the previous absolute time. The difference between this time and the current absolute time determines the delay before this event should be repeated. The event is then added to the queue at the appropriate position for the event's next occurrence. This strategy insures that even if some delay occurs in dispatching the event, following events occur at the proper time. After all events with a offset of zero have been processed, the SC queues a new timer event at the offset to the first waiting event, after which the above process repeats.

The decoupling of the absolute clock value and the actual dispatching of clock events is critical in implementing some of the unique and novel aspects of the invention, *e.g.* the ability to control easily the rate at which playback of the stream occurs. When initially created, the SC records a value referred to as the rate multiplier, which is initially set to one. Whenever the SC requests the underlying operating system to queue a timer event, the actual time delay requested is multiplied by the rate multiplier. A rate multiplier greater than one results in faster playback of the stream, while a multiplier of less than one results in a slower playback of the stream.

For the LCs to position themselves properly (see below), the same rate multiplier must be used to modify the Presentation Time Stamp (PTS) stored in each block when calculating positioning. Thus, the SC makes a small set of functions available to the LC which perform comparisons and operations on PTS values, after properly synchronizing those values with the rate multiplier stored in the SC.

The combination of the Stream Clock, with independent event generation for each LC and the ability to modify logically the rate at which all events are dispatched using a single value provides a novel and unique application of common time-based software design techniques.

THE LINEAR CACHE

A key component of the invention is the Linear Cache (LC). The LC is used to store a single stream of digital information in proper sequence. The LC maintains this stream as a list of the digital blocks that make up the stream, and assumes that each block has a unique size and attributes. Each block within the stream is marked with a Presentation Time Stamp (PTS) which indicates when that block should be presented to the decoding process, be it hardware or software. The PTS is a monotonically increasing value initialized to zero when the LC first begins operation on a stream, and is not related to any time stamps defined by the underlying stream encoding technique except that once the LC begins operation, there is a constant offset between the PTS stored in the LC and any time stamps stored within the encoded digital blocks. It is assumed that the decoding process properly handles embedded time stamps. Figure 3 gives a pictorial overview of the preferred embodiment of the Linear Cache.

The capture mechanism for a particular stream type gives each encoded digital block to the LC as it arrives (301). The LC marks that block with the current PTS for the stream. The LC maintains a window (302), which is defined as the maximum time, based on PTS values, between when the newest block in the LC arrived and the oldest block which the LC may hold. Thus, the window represents a time span into the past history of the stream. The LC discards blocks which fall outside the window (303), thus the window allows one to look a fixed distance into the past history of a stream. This mechanism allows choices and tradeoffs between available storage space and the availability of past information for viewing.

The capture mechanism is responsible for providing certain attributes along with each new block (304). The first of these is the key frame attribute, which indicates that this block begins a sequence of interrelated blocks. When performing random access operations on a stream, the LC only allows positioning of the stream to a block marked as a key frame. The second is the End Of Segment (EOS) attribute, which indicates that the stream has ended, and no more data are to be presented. For example, the I-frame of an MPEG GOP is marked as a key frame by the capture mechanism, but all other blocks are not so marked.

The LC maintains an indication of the next block to be presented to the decoding process, which is referred to as the current block indicator (305). For normal viewing, this block is the block last added to the LC by the capture mechanism. Thus, the stream is presented live, with at most a one-frame time delay between capture and presentation.

Random access to the information stream is achieved by moving the current block indicator to some other block in the LC. If the LC is requested to move the indicator, and the current block indicated is not a key frame, the LC instructs the decoding process to reset its decoding state, thus purging any partially constructed presentation data. For example, if the current frame is an MPEG predictive frame (P-frame), then the decoding process has state information, including the current picture to be modified by the P-frame. This data must be purged so that the decoder begins in the proper state.

Referring again to Figure 2, the BC requests repositioning of an LC by specifying a PTS value. The LC finds the block containing a PTS which is closest to that requested by the BC. There are two unique cases for repositioning: the first, if the LC is marked as the key stream, and the second when it is not.

If the LC is marked as a key stream, the LC scans the blocks in the cache to find the key frame which is nearest to the requested PTS value, searching both before and after the desired value. Once properly positioned, the LC returns to the BC the PTS of the key frame block which was identified.

Following positioning of the key stream, the BC instructs each remaining stream to position itself to the PTS returned by the key-stream. Key frames in other LCs may not align with those in the key stream. Each LC handles this problem by requesting the decoding process to purge any decoding state, and then stores an indication that suppresses the LC from actually presenting blocks to the decoding process.

Following this, the LC behaves normally, accepting clock events and advancing the current block indicator to match, except that the blocks are not actually presented to the decoding process. When the LC encounters a key frame, it deletes the indication suppressing presentation of blocks, and presents the current block to the decoding process. This brings the stream into full synchronization with the key stream.

The following subsections describe in detail how positioning operations are handled by the LC.

Forward

The forward function is implemented by moving the current block indicator forward through the cache by one block for each event generated by the Stream Clock. Each successive block in the cache is given to the decoding process for presentation as clock events are dispatched by the SC. A rate multiplier of one causes the current block indicator to move forward one block in the LC at the same rate as new blocks are added to the LC by the capture process. Conceptually, the current block indicator points to the same position in the cache, while the window of cached blocks moves beneath it. From the perspective of the consumer of a stream, this results in play out at a fixed delay from real time.

There exists a rate at which the decoding process may not be able to present a block to the consumer before decoding of the next block must begin, *i.e.* the decoder can not keep up with the playback rate requested. The LC recognizes this case, and reverts to a key frame based method, where intermediate blocks are skipped and only key frames are presented to the decoder.

The current block indicator can not be positioned to an uncached block. If the rate multiplier indicates a rate greater than one, then the current block indicator eventually moves to the front of the cache. In this case, the key frame LC indicates to the BC that no more blocks are present in the cache. The BC resets the SC rate multiplier to one, and each LC positions the current block indicator to the latest cached block. The effect of this operation is to revert to undelayed play out of the actual information stream.

Reverse

The reverse function is implemented by moving the current block indicator backwards through the cache by one block for each clock event generated by the SC. Again, the key stream LC is used to control positioning. The BC instructs each LC to move to reverse mode, in which the current block indicator is moved backwards one block on each clock event. In addition, this

mode causes the LC to suppress¹³ presentation of all but key frames to the decoding process. When the current block indicator is moved over a key frame, the block is presented to the decoding process.

The rate at which blocks are presented to the decoding process is controlled by the rate multiplier in the Stream Clock, allowing for arbitrary speed of reverse operation.

The current block indicator can not be moved past the oldest block in the cache. If a clock event would result in moving the current block indicator past the earliest block, then the indicator is set to that block, and play out continues as described above. The key frame LC indicates to the BC that the oldest cached block was reached. The BC resets the SC rate multiplier to one, and each LC positions the current block indicator to the oldest cached block. Externally, it appears as if the stream began playing in a forward direction again, time delayed by the window size.

Pause

The BC implements the pause function by locking the current block indicator in the key stream LC to that block. The LC can only lock the indicator to a key frame block, thus it searches forward for such a block in the cache. If no key frame is present, the LC stores an indication that a lock has been requested. When the capture process presents a key frame, the LC thus recognizes that the lock was requested, and locks the current block indicator to that key frame. The LC also presents this key frame to the decoding process, such that proper positioning, from the decoding process point of view, is maintained. Following this, as new blocks are added to the front of the cache, the block indicated moves backwards in the cache. Additionally, the LC sets the indicator suppressing further presentation of data to the decoding process.

For each non-key-stream LC, the BC simply suppresses further presentation of data blocks to the decoding process. The BC may be requested to release the pause (*i.e.* to play the stream). In this case, the BC obtains the PTS of the current block (as modified by the SC rate multiplier) from the key stream LC, and directs each remaining LC to position appropriately as described above. It then releases the lock on each LC, resulting in presentation of blocks to the decoding processes.

It is possible that the window moves past the block to which the indicator is locked because capture of the incoming stream is continuous. In this case, the indicator is unlocked and the left pointing at the earliest block in the cache. This gives the external effect of automatic release of the pause, similar to the behavior of a VCR, which releases the pause to avoid excessive wear on the tape and playback heads.

Get Clip

This function is used to extract a range of blocks from the cache. There are two modes in which it operates: relative mode and absolute mode.

As discussed above, the cached blocks appear as a linear list of blocks in presentation order. The `getclip()` function returns a pointer to a cached block, plus a count of the number of blocks that make up the range. This information is referred to as a clip. For efficiency sake, the pointer returned directly indicates a cached block, and the following blocks are retrieved by following the next block link in each cached block. This list is only valid if the first block pointed to and all succeeding blocks remain cached until the caller has completed its operation on the clip.

The LC has the following behavior to avoid invalidating blocks which are part of a clip: the `getclip()` function sets an indicator in the LC that suppresses discarding blocks which fall outside the window. This indicator is cleared when the `releaseclip()` function is called. The `releaseclip()` function also trims the cache to the window size.

The result of this sequence is that the number of blocks in the LC may temporarily increase. It is the responsibility of the caller of the `getclip()` function to operate on the clip as rapidly as possible because storage space for cached blocks is a limited resource. However, this strategy allows a clip to be captured without interfering with the ongoing playback of blocks from the LC, which avoids presenting the consumer of the blocks with any artifacts of the capture operation.

In relative mode, the clip to be captured is specified as two numbers, the first representing an offset from the current block indicator into previous blocks in the cache, and the second representing an offset from the current block indicator into subsequent blocks in the cache. The offset in both cases is

expressed as a time differential from the PTS stored in the current block. For instance, a capture operation might occur from five seconds before the current block to ten seconds after the current block. No uncached blocks may be indicated by the `getclip()` operation. Thus, the LC trims the actual clip returned to match the blocks available in the cache.

In absolute mode, the clip to be captured is specified as two numbers, the first representing an absolute time stamp based on when the LC began operation, and the second representing a time offset from that value, which may be a negative offset (indicating a distance backwards in the cache) or a positive value (indicating distance forwards in the cache). The clip is returned as described above, and trimmed as necessary. This also implies that `getclip()` can return an empty list of blocks if no block in the cache is within the time window described.

CLIP CAPTURE

The Clip Capture module is responsible for selecting ranges of blocks from a set of LCs, bringing them together into an appropriate storage or transmission format, and saving or sending them as needed. This collection of blocks from multiple LCs is also referred to as a clip. The Clip Capture operation is invoked by the BC on external request.

There are two ways in which this request may be made:

First, the BC may indicate that capture should occur based on the current block indicator in the key stream LC, which is referred to as a relative capture, in which case a relative range of capture is specified.

Second, the BC may indicate an absolute range, in which case only blocks marked with a PTS inclusive in that range are captured. This is referred to as an absolute capture.

The operation of the Clip Capture module is straightforward. For each LC, the CC module calls the `getclip()` function of that LC with parameters as directed by the BC. The Clip Capture module then linearizes the blocks into a multiplexed stream appropriate for the final clip format desired. For example, it might generate an MPEG System Stream, interleaving blocks from the various LCs as needed.

This interleaving is conceptually simple. The CC loops through each clip obtained from an LC, choosing the block with the earliest PTS from among all clips, outputting the block as appropriate, and advancing the block pointer for that clip to the next block. When all the blocks from all clips are exhausted, the capture operation is complete. The CC then loops through all LCs calling the releaseclip() function.

Although the invention is described herein with reference to the preferred embodiment, one skilled in the art will readily appreciate that other applications may be substituted for those set forth herein without departing from the spirit and scope of the present invention. Accordingly, the invention should only be limited by the Claims included below.

17
CLAIMS

1. A method for providing pass through or capture of continuous linear streams of digital information represented in various formats while providing the appearance of a locally stored stream, comprising the steps of:

providing at least one media cache for copying blocks of data from said information stream, wherein data in said media cache can be viewed as a snapshot of said continuous stream of digital information;

providing a playback pointer, wherein a playback pointer position selects a portion of said media cache that is to be accessed to provide functions including any of pause, rewind, fast forward, play, play faster, play slower, and play in reverse; and

wherein said underlying information stream can be composed of a plurality of discrete blocks of data.

2. The method of Claim 1, further comprising the step of:

providing a buffer controller that accepts external requests for operations on said information stream and that, in turn, generates appropriate control messages; and

wherein said buffer controller instantiates methods that provide for system operation in connection with any particular stream of information.

3. The method of Claim 1, further comprising the step of:

providing a stream clock for synchronizing operations on said more than one media cache, wherein multiple streams of data which must be presented in a synchronized fashion are correctly positioned, and wherein said multiple streams of data present their data at consistent delivery rates; and

wherein said stream clock provides a central synchronization facility that distributes time based events to a number of media caches.

4. The method of Claim 3, further comprising the step of:

designating one of said multiple streams as a key stream;

wherein all positioning of playback from each media cache is based on first positioning said key stream, and then using a final position of that stream to position all other streams properly;

wherein positioning is based on aligning a current block pointer to the beginning of an MPEG Group Of Pictures when said key stream is an MPEG video stream; and

wherein a buffer controller¹⁸ repositions all other streams to be synchronized with said key stream once said key stream is positioned.

5. The method of Claim 3, further comprising the step of:
providing a buffer controller that is responsible for controlling a clock rate of said stream clock;

wherein the rate at which said stream clock delivers clock events to said media cache is based on a rate specified by said buffer controller; and

wherein said buffer controller can speed up or slow down playback of an entire collection of streams by requesting said stream clock to adjust its clock rate.

6. The method of Claim 3, wherein events are queued with a relative offset from a previously queued event;

wherein a first event in said queue is marked with a time which is an offset from a current time;

wherein said stream clock requests that an underlying operating system invoke a call out when that amount of time has passed; and

wherein said stream clock dispatches a first clock event and then each succeeding event whose offset is zero when said call out is invoked.

7. The method of Claim 1, wherein said media cache maintains said stream as a list of the digital blocks that make up said stream, and assumes that each block in said stream has a unique size and attributes;

wherein said attributes include any of a key frame attribute which indicates that said block begins a sequence of inter-related blocks such that said media cache only allows positioning of said stream to a block marked as a key frame when performing random access operations on said stream; and an End Of Segment (EOS) attribute which indicates that said stream has ended and that no more data are to be presented;

wherein each block within said stream is marked with a Presentation Time Stamp (PTS) which indicates when that block should be presented to a decoding process;

wherein said PTS is a monotonically increasing value initialized to zero when said media cache first begins operation on a stream and is not related to any time stamps defined by an underlying stream encoding technique, except that there is a constant offset between said PTS stored in said media cache and any time stamps stored within said encoded digital blocks once said media cache begins operation;

wherein a capture mechanism for a particular stream type gives each encoded digital block to said media cache as it arrives and provides certain attributes along with each new block;

wherein said media cache marks said encoded digital block with a current PTS for said stream;

wherein said media cache maintains a window which is the maximum time, based on PTS values, between the time when a newest block in said media cache has arrived and an oldest block which said media cache may hold, such that said window represents a time-span into a past history of said stream;

wherein said media cache maintains an indication of a next block to be presented to a decoding process; and

wherein said next block is a block last added to said media cache by said capture mechanism, such that said stream is presented live with at most a one-frame time delay between capture and presentation.

8. The method of Claim 7, wherein random access to said stream is achieved by moving a current block indicator to another block in said media cache; and

wherein said media cache instructs a decoding process to reset its decoding state, thus purging any partially constructed presentation data, if said media cache is requested to move said block indicator, and a current block indicated is not a key frame.

9. The method of Claim 7, further comprising the step of:

providing a buffer controller that requests repositioning of a media cache by specifying a PTS value;

wherein said media cache scans blocks of data in said media cache to find a key frame which is nearest to a requested PTS value, searching both before and after a desired value, if said media cache is marked as a key stream, wherein said media cache returns PTS of said key frame block which was identified to said buffer controller, once said media cache is properly positioned;

instructing each remaining stream to position itself to said PTS returned by said key stream, wherein key frames in other media caches may not align with those in said key stream, in which event each media cache requests a decoding process to purge any decoding state, and then stores an indication that suppresses each other media cache from actually presenting blocks to said decoding process;

wherein said media cache accepts clock events and advances a current block indicator to match, except that said blocks are not actually presented to said decoding process;

deleting an indication suppressing presentation of blocks; and

presenting a current block to said decoding process when said media cache encounters a key frame.

10. The method of Claim 1, wherein a forward/reverse function is implemented by moving a current block indicator forward/reverse through said media cache by one block for each event generated by a stream clock;

wherein each successive block in said media cache is given to an optional decoding process for presentation as clock events are dispatched by said stream clock;

wherein a rate multiplier causes a current block indicator to move forward/reverse within said media cache at a selected rate as new blocks are added to said media cache by a capture process, resulting in play out at a fixed delay from real time;

wherein said media cache reverts to a key frame based method in which intermediate blocks are skipped and only key frames are presented to said decoder if said decoder can not keep up with a playback rate requested;

wherein a buffer controller resets said stream clock rate multiplier to one, and each media cache positions a current block indicator to a latest cached block if said current block indicator moves to a front of said media cache where no more blocks are present in said media cache;

wherein said current block indicator is set to an earliest block in said media cache if a clock event would result in moving a current block indicator past said earliest block; and

wherein play out thereafter continues and that it appears as if said stream is playing in a forward direction again.

11. The method of Claim 1, further comprising the step of:

providing a buffer controller for implementing a pause function by locking a current block indicator in a key stream media cache to a key frame block nearest a currently indicated block;

searching forward for said key frame block in said media cache if said current block is not a key frame block;

storing an indication that a lock has been requested if no key frame is present;

locking said current block indicator to a key frame when a capture process presents a key frame; and

wherein said indicator is unlocked and left pointing at an earliest block in said media cache if a stream window moves past a block to which said indicator is locked, thereby providing automatic release of said pause.

12. The method of Claim 1, further comprising the steps of;

extracting a range of blocks from said media cache in either of a relative mode and an absolute mode by returning a pointer to a cached block plus a count of the number of blocks that make up an extract range; and

providing an indicator in said media cache that suppresses discarding blocks within said extraction which fall outside a window.

13. The method of Claim 1, further comprising the step of:

providing a clip capture module for selecting ranges of blocks from a set of media caches, bringing said blocks together into an appropriate storage or transmission format, and saving or sending said blocks as needed; and

wherein a buffer controller indicates where a capture should occur based on a current block indicator in a key stream media cache.

14. An apparatus for providing pass through or capture of continuous linear streams of digital information represented in various formats while providing the appearance of a locally stored stream, comprising:

at least one media cache for copying blocks of data from said information stream, wherein data in said media cache can be viewed as a snapshot of said continuous stream of digital information;

a playback pointer, wherein a playback pointer position selects a portion of said media cache that is to be accessed to provide functions including any of pause, rewind, fast forward, play, play faster, play slower, and play in reverse; and

wherein said underlying information stream can be composed of a plurality of discrete blocks of data.

15. The apparatus of Claim 14, further comprising:

a buffer controller that accepts external requests for operations on said information stream and that, in turn, generates appropriate control messages; and

wherein said buffer controller instantiates methods that provide for system operation in connection with any particular stream of information.

16. The apparatus of Claim 14, further comprising:
a stream clock for synchronizing operations on said more than one media cache, wherein multiple streams of data which must be presented in a synchronized fashion are correctly positioned, and wherein said multiple streams of data present their data at consistent delivery rates; and
wherein said stream clock provides a central synchronization facility that distributes time based events to a number of media caches.
17. The apparatus of Claim 16, further comprising:
a module for designating one of said multiple streams as a key stream;
wherein all positioning of playback from each media cache is based on first positioning said key stream, and then using a final position of that stream to position all other streams properly;
wherein positioning is based on aligning a current block pointer to the beginning of an MPEG Group Of Pictures when said key stream is an MPEG video stream; and
wherein a buffer controller repositions all other streams to be synchronized with said key stream once said key stream is positioned.
18. The apparatus of Claim 16, further comprising:
a buffer controller that is responsible for controlling a clock rate of said stream clock;
wherein the rate at which said stream clock delivers clock events to said media cache is based on a rate specified by said buffer controller; and
wherein said buffer controller can speed up or slow down playback of an entire collection of streams by requesting said stream clock to adjust its clock rate.
19. The apparatus of Claim 16, wherein events are queued with a relative offset from a previously queued event;
wherein a first event in said queue is marked with a time which is an offset from a current time;
wherein said stream clock requests that an underlying operating system invoke a call out when that amount of time has passed; and
wherein said stream clock dispatches a first clock event and then each succeeding event whose offset is zero when said call out is invoked.

20. The apparatus of Claim 14,²³ wherein said media cache maintains said stream as a list of the digital blocks that make up said stream, and assumes that each block in said stream has a unique size and attributes;

wherein said attributes include any of a key frame attribute which indicates that said block begins a sequence of inter-related blocks such that said media cache only allows positioning of said stream to a block marked as a key frame when performing random access operations on said stream; and an End Of Segment (EOS) attribute which indicates that said stream has ended and that no more data are to be presented;

wherein each block within said stream is marked with a Presentation Time Stamp (PTS) which indicates when that block should be presented to a decoding process;

wherein said PTS is a monotonically increasing value initialized to zero when said media cache first begins operation on a stream and is not related to any time stamps defined by an underlying stream encoding technique, except that there is a constant offset between said PTS stored in said media cache and any time stamps stored within said encoded digital blocks once said media cache begins operation;

wherein a capture mechanism for a particular stream type gives each encoded digital block to said media cache as it arrives and provides certain attributes along with each new block;

wherein said media cache marks said encoded digital block with a current PTS for said stream;

wherein said media cache maintains a window which is the maximum time, based on PTS values, between the time when a newest block in said media cache has arrived and an oldest block which said media cache may hold, such that said window represents a time-span into a past history of said stream;

wherein said media cache maintains an indication of a next block to be presented to a decoding process; and

wherein said next block is a block last added to said media cache by said capture mechanism, such that said stream is presented live with at most a one-frame time delay between capture and presentation.

21. The apparatus of Claim 20, wherein random access to said stream is achieved by moving a current block indicator to another block in said media cache; and

wherein said media cache instructs a decoding process to reset its decoding state, thus purging any partially constructed presentation data, if said

media cache is requested to move said block indicator, and a current block indicated is not a key frame.

22. The apparatus of Claim 20, further comprising:

- a buffer controller that requests repositioning of a media cache by specifying a PTS value;

- wherein said media cache scans blocks of data in said media cache to find a key frame which is nearest to a requested PTS value, searching both before and after a desired value, if said media cache is marked as a key stream, wherein said media cache returns PTS of said key frame block which was identified to said buffer controller, once said media cache is properly positioned;

- a module for instructing each remaining stream to position itself to said PTS returned by said key stream, wherein key frames in other media caches may not align with those in said key stream, in which event each media cache requests a decoding process to purge any decoding state, and then stores an indication that suppresses each other media cache from actually presenting blocks to said decoding process;

- wherein said media cache accepts clock events and advances a current block indicator to match, except that said blocks are not actually presented to said decoding process;

- a module for deleting an indication suppressing presentation of blocks;
- and

- a module for presenting a current block to said decoding process when said media cache encounters a key frame.

23. The apparatus of Claim 14, wherein a forward/reverse function is implemented by moving a current block indicator forward/reverse through said media cache by one block for each event generated by a stream clock;

- wherein each successive block in said media cache is given to an optional decoding process for presentation as clock events are dispatched by said stream clock;

- wherein a rate multiplier causes a current block indicator to move forward/reverse within said media cache at a selected rate as new blocks are added to said media cache by a capture process, resulting in play out at a fixed delay from real time;

- wherein said media cache reverts to a key frame based method in which intermediate blocks are skipped and only key frames are presented to said decoder if said decoder can not keep up with a playback rate requested;

wherein a buffer controller resets said stream clock rate multiplier to one, and each media cache positions a current block indicator to a latest cached block if said current block indicator moves to a front of said media cache where no more blocks are present in said media cache;

wherein said current block indicator is set to an earliest block in said media cache if a clock event would result in moving a current block indicator past said earliest block; and

wherein play out thereafter continues and that it appears as if said stream is playing in a forward direction again.

24. The apparatus of Claim 14, further comprising:

a buffer controller for implementing a pause function by locking a current block indicator in a key stream media cache to a key frame block nearest a currently indicated block;

a module for searching forward for said key frame block in said media cache if said current block is not a key frame block;

a module for storing an indication that a lock has been requested if no key frame is present;

a module for locking said current block indicator to a key frame when a capture process presents a key frame; and

wherein said indicator is unlocked and left pointing at an earliest block in said media cache if a stream window moves past a block to which said indicator is locked, thereby providing automatic release of said pause.

25. The apparatus of Claim 14, further comprising:

a module for extracting a range of blocks from said media cache in either of a relative mode and an absolute mode by returning a pointer to a cached block plus a count of the number of blocks that make up an extract range; and

an indicator in said media cache that suppresses discarding blocks within said extraction which fall outside a window.

26. The apparatus of Claim 14, further comprising:

a clip capture module for selecting ranges of blocks from a set of media caches, bringing said blocks together into an appropriate storage or transmission format, and saving or sending said blocks as needed; and

wherein a buffer controller indicates where a capture should occur based on a current block indicator in a key stream media cache.

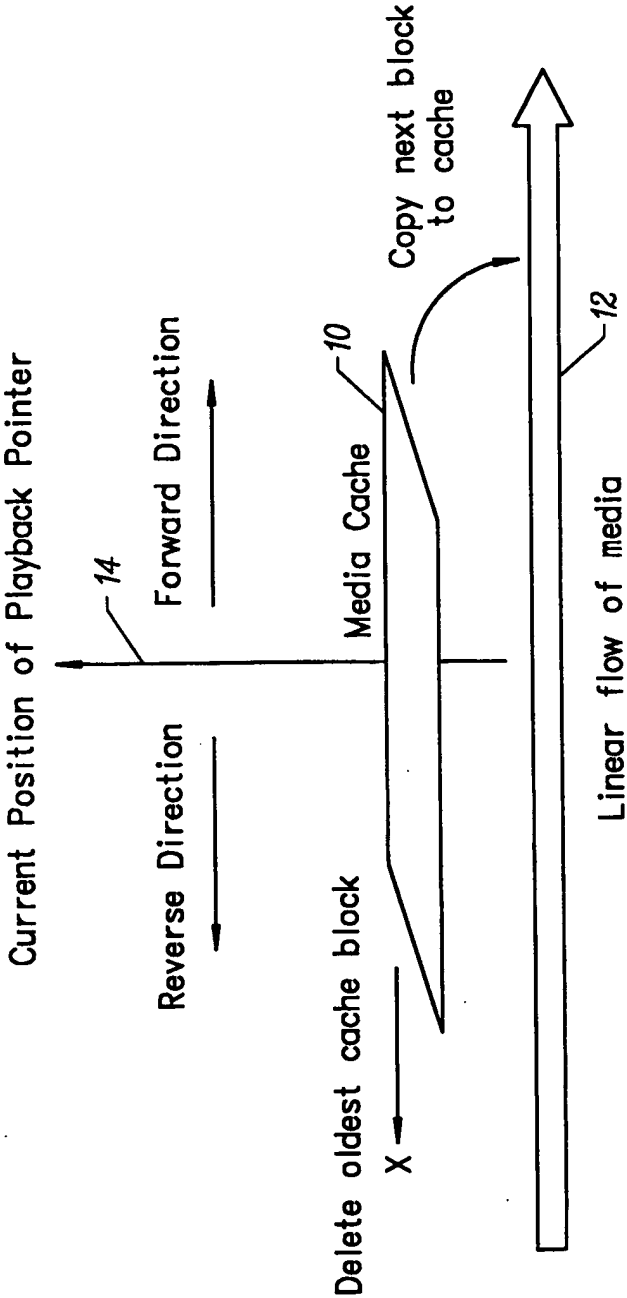


FIG. 1

2/3

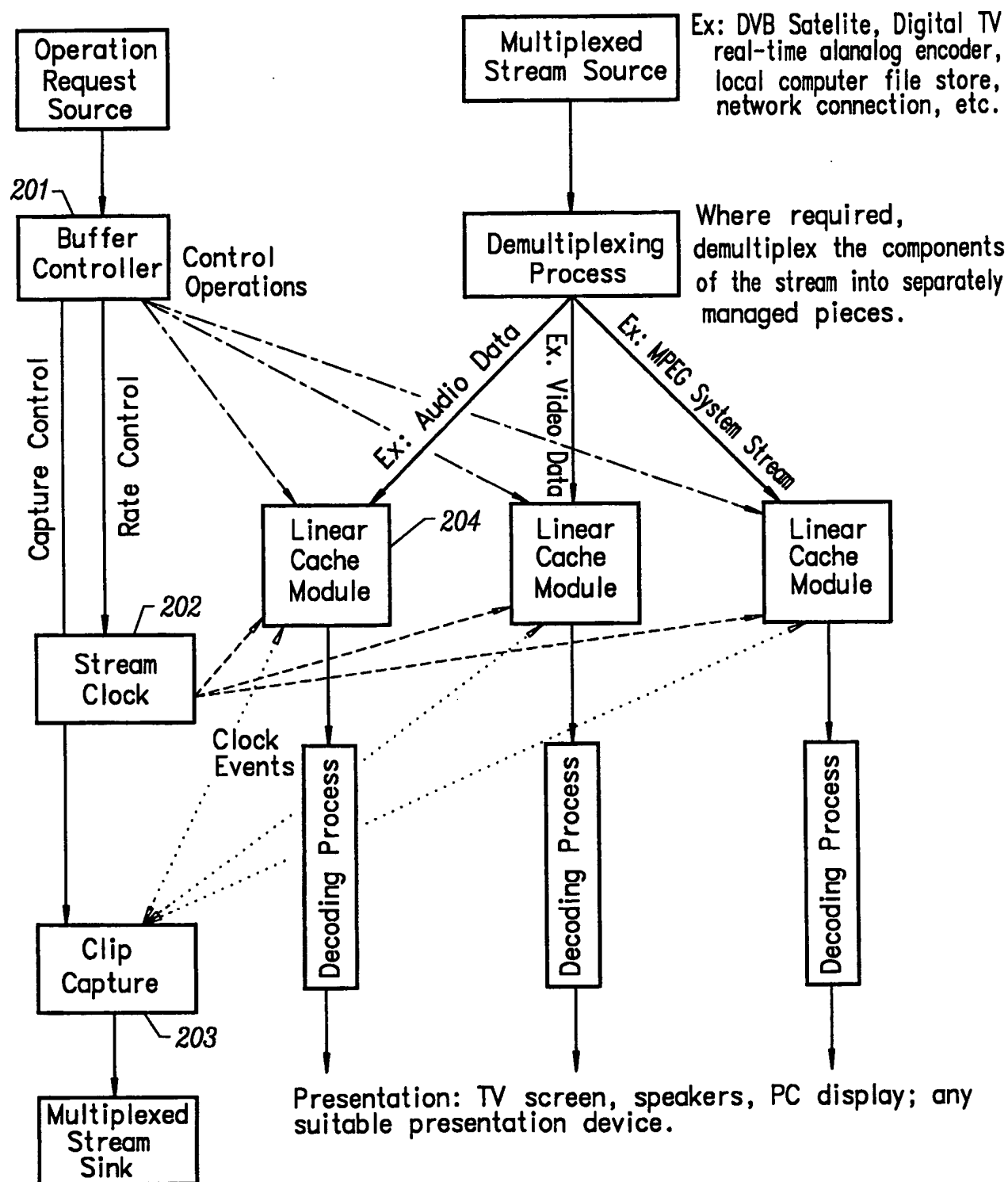


FIG. 2

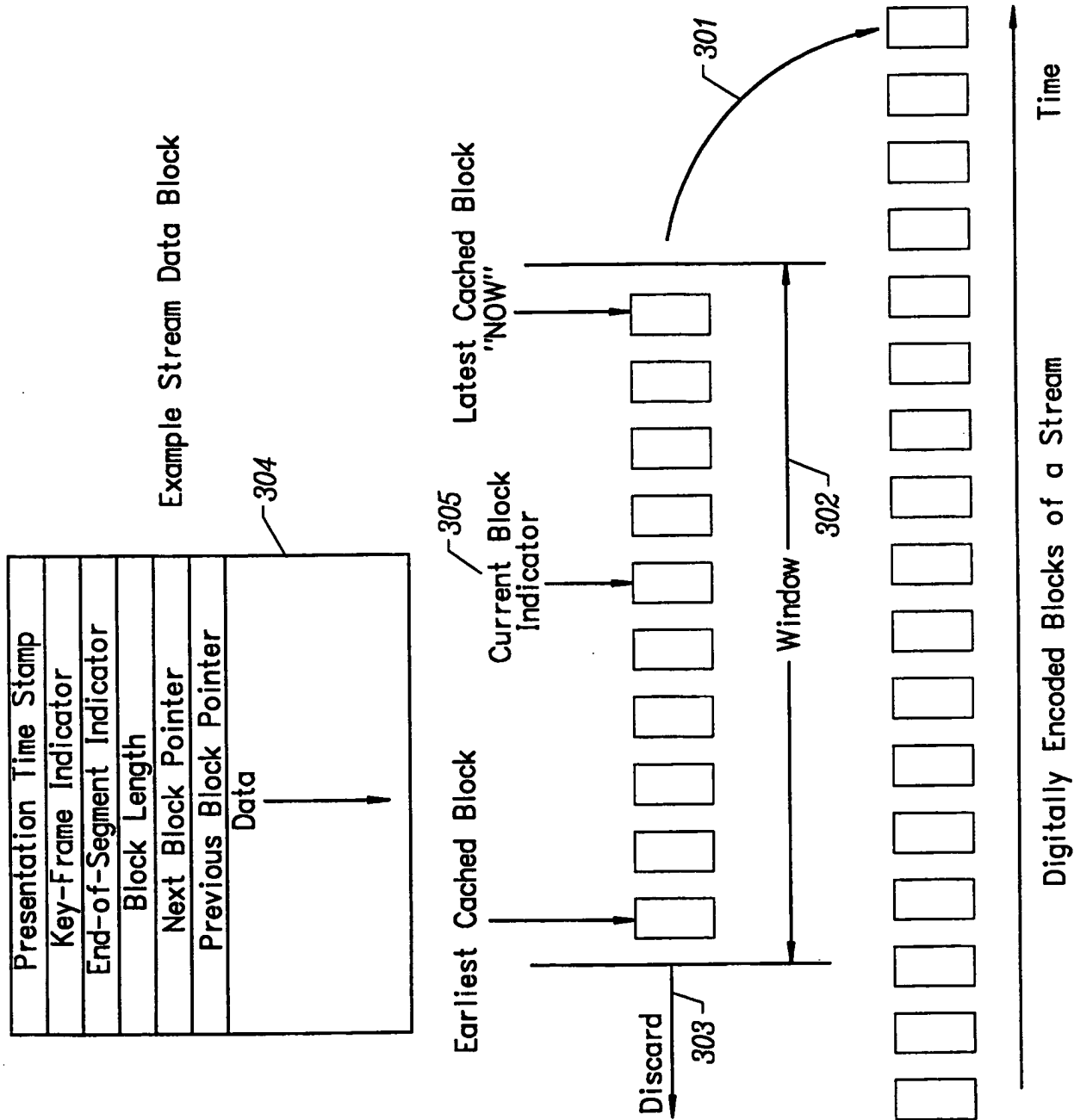


FIG. 3

A. CLASSIFICATION OF SUBJECT MATTER

IPC 6 H04N5/775

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 H04N

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	US 5 438 423 A (LYNCH ET AL.) 1 August 1995 see the whole document ---	1, 14
A	EP 0 726 574 A (MATSUSHITA ELECTRIC INDUSTRIAL CO., LTD) 14 August 1996 see column 8, line 56 - column 16, line 17; figures 1-6 ---	1, 14
A	US 5 706 388 A (ISAKA) 6 January 1998 see column 2, line 29 - column 6, line 23; figures 1-3 ---	1, 14
A	US 5 696 868 A (KIM ET AL.) 9 December 1997 see the whole document -----	1, 14



Further documents are listed in the continuation of box C.



Patent family members are listed in annex.

* Special categories of cited documents :

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"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

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"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

"Z" document member of the same patent family

Date of the actual completion of the international search

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01/07/1999

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